

# Module 5: Dictionary ADT (Hash table)

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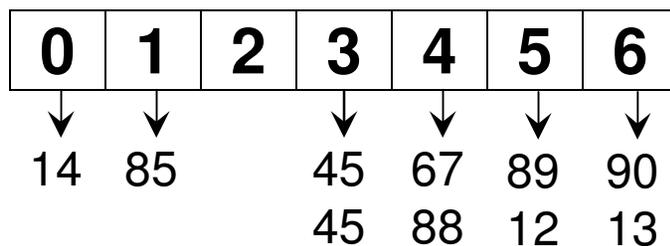
# Dictionary ADT

- Models data as a collection of key-value pairs.
  - The Keys are unique
  - A value has a unique key and is accessed in the dictionary using that key
- Operations
  - Insert: The addition of a key-value pair
  - Delete: The removal of a key-value pair
  - Find: Lookup the existence of a key-value pair
  - isEmpty: Whether the dictionary is empty
- Implementations: Hash tables, Binary search trees

# Hash table

- Maps the elements (values) of a collection to a unique key and stores them as key-value pairs.
- Hash table of size  $m$  (where  $m$  is the number of unique keys, ranging from 0 to  $m-1$ ) uses a hash function  $H(v) = v \bmod m$
- The hash value (a.k.a. hash index) for an element  $v$  is  $H(v) = v \bmod m$  and corresponds to one of the keys of the hash table.
- The size of the Hash table is typically a prime integer.
- Example: Consider a hash table of size 7. Its hash function is  $H(v) = v \bmod 7$ .
- Let an array  $A = \{45, 67, 89, 45, 85, 12, 88, 90, 13, 14\}$

|                    |    |    |    |    |    |    |    |    |    |    |
|--------------------|----|----|----|----|----|----|----|----|----|----|
| Value, $v$         | 45 | 67 | 89 | 45 | 85 | 12 | 88 | 90 | 13 | 14 |
| $H(v) = v \bmod 7$ | 3  | 4  | 5  | 3  | 1  | 5  | 4  | 6  | 6  | 0  |



**We will implement Hash table as an array of singly linked lists**

# Space-Time Tradeoff

- Note: At the worst case, there could be only one linked list in the hash table (i.e., all the elements map to the same key).
- On average, we expect the 'n' elements to be evenly divided across the 'm' keys, so that the length of a linked list is  $n/m$ . Nevertheless, for a hash table of certain size (m), 'n' is the only variable.
- Space complexity:  $\Theta(n)$ 
  - For an array of 'n' elements, we need to allocate space for 'n' nodes (plus the 'm' head node) across the 'm' linked lists.
  - Since usually,  $n \gg m$ , we just consider the overhead associated with storing the 'n' nodes
- Time complexity:
  - Insert/Delete/Lookup:  $O(n)$ , we may have to traverse the entire linked list
  - isEmpty:  $O(m)$ , we have to check whether each index in the Hash table has an empty linked list or not.

# Example: Number of Comparisons

Array, A = {45, 23, 11, 78, 89, 44, 22, 28, 41, 30}

$H(v) = v \text{ mod } 7$

| Hash table | 0  | 1  | 2  | 3  | 4  | 5  | 6  | Successful Search, # comparisons |
|------------|----|----|----|----|----|----|----|----------------------------------|
|            | ↓  | ↓  | ↓  | ↓  | ↓  | ↓  | ↓  |                                  |
|            | 28 | 78 | 23 | 45 | 11 | 89 | 41 | 1                                |
|            |    | 22 | 44 |    |    |    |    | 2                                |
|            |    |    | 30 |    |    |    |    | 3                                |

**Average Number of Comparisons for a Successful Search (Hash table)**

$$= \frac{(7 \cdot 1) + (2 \cdot 2) + (1 \cdot 3)}{10} = \frac{14}{10} = 1.4$$

**Worst Case Number of Comparisons for a Successful Search (Hash table) = 3**

**Worst Case Number of Comparisons for an Unsuccessful Search (Hash table) = 3**

# Example: Number of Comparisons

Array, A = {45, 23, 11, 78, 89, 44, 22, 28, 41, 30}

$H(v) = v \text{ mod } 7$

| Hash table | 0  | 1  | 2  | 3  | 4  | 5  | 6  | Successful Search, # comparisons |
|------------|----|----|----|----|----|----|----|----------------------------------|
|            | ↓  | ↓  | ↓  | ↓  | ↓  | ↓  | ↓  | 1                                |
|            | 28 | 78 | 23 | 45 | 11 | 89 | 41 | 2                                |
|            |    | 22 | 44 |    |    |    |    | 3                                |
|            |    |    | 30 |    |    |    |    |                                  |

|  |  |
|--|--|
| <p><b>Average Number of Comparisons for a Successful Search (Hash table)</b></p> | $= \frac{(7 \cdot 1) + (2 \cdot 2) + (1 \cdot 3)}{10} = \frac{14}{10} = 1.4$ |
|--|--|

|   |  |
|---|--|
| <p><b>Average Number of Comparisons for a Successful Search (Array)</b></p> | $= \frac{1 + 2 + 3 + \dots + 10}{10} = \frac{10 \cdot 11 / 2}{10} = 5.5$ |
|---|--|

|  |  |                                      |
|--|--|--------------------------------------|
| <p><b>Worst Case Number of Comparisons For a Successful Search</b></p> | <p><b>Hash table</b></p> <p><b>3</b></p> | <p><b>Array</b></p> <p><b>10</b></p> |
| <p><b>For an unsuccessful Search</b></p>                               | <p><b>3</b></p>                          | <p><b>10</b></p>                     |

# Load Imbalance Index

- Load Imbalance Index is a measure of the efficiency in using the memory allocated for a hash table.
  - The larger the value for the index, the larger the imbalance and vice-versa
  - The index ranges from 0 to 1.

$$\frac{L_{\max} - L_{\min}}{L_{\max} + L_{\min}}$$

## Example 1:

|          |          |          |          |          |          |          |
|----------|----------|----------|----------|----------|----------|----------|
| <b>0</b> | <b>1</b> | <b>2</b> | <b>3</b> | <b>4</b> | <b>5</b> | <b>6</b> |
| ↓        | ↓        | ↓        | ↓        | ↓        | ↓        | ↓        |
| 28       | 78       | 23       | 45       | 11       | 89       | 41       |
|          | 22       | 44       |          |          |          |          |
|          |          | 30       |          |          |          |          |

### Load Imbalance Index

$$= \frac{(3 - 1)}{(3 + 1)} = \frac{2}{4} = 0.50$$

## Example 2:

|          |          |          |          |          |          |          |          |          |          |           |
|----------|----------|----------|----------|----------|----------|----------|----------|----------|----------|-----------|
| <b>0</b> | <b>1</b> | <b>2</b> | <b>3</b> | <b>4</b> | <b>5</b> | <b>6</b> | <b>7</b> | <b>8</b> | <b>9</b> | <b>10</b> |
| ↓        | ↓        | ↓        | ↓        | ↓        | ↓        | ↓        | ↓        | ↓        | ↓        | ↓         |
| 22       | 78       | -        | -        | -        | -        | 28       | -        | 30       | -        | -         |
| 44       | 23       |          |          |          |          |          |          | 41       |          |           |
| 11       | 45       |          |          |          |          |          |          |          |          |           |
|          | 89       |          |          |          |          |          |          |          |          |           |

### Load Imbalance Index

$$= \frac{(4 - 0)}{(4 + 0)} = \frac{4}{4} = 1.0$$

# Review of Singly Linked List

```
private:
    Node *headPtr;

public:
    List(){
        headPtr = new Node();
        headPtr->setNextNodePtr(0);
    }

    Node* getHeadPtr(){
        return headPtr;
    }

    bool isEmpty(){
        if (headPtr->getNextNodePtr() == 0)
            return true;

        return false;
    }
}
```

**Class List**  
**C++**

```
void insert(int data){
    Node* currentNodePtr = headPtr->getNextNodePtr();
    Node* prevNodePtr = headPtr;

    while (currentNodePtr != 0){
        prevNodePtr = currentNodePtr;
        currentNodePtr = currentNodePtr->getNextNodePtr();
    }

    Node* newNodePtr = new Node();
    newNodePtr->setData(data);
    newNodePtr->setNextNodePtr(0);
    prevNodePtr->setNextNodePtr(newNodePtr);
}
```

```
bool deleteElement(int deleteData){  
    Node* currentNodePtr = headPtr->getNextNodePtr();  
    Node* prevNodePtr = headPtr;  
    Node* nextNodePtr = headPtr;  
  
    while (currentNodePtr != 0){  
        if (currentNodePtr->getData() == deleteData){  
            nextNodePtr = currentNodePtr->getNextNodePtr();  
            prevNodePtr->setNextNodePtr(nextNodePtr);  
            return true;  
        }  
  
        prevNodePtr = currentNodePtr;  
        currentNodePtr = currentNodePtr->getNextNodePtr();  
    }  
    return false;  
}
```

**deleteElement(int)**

**C++**

```
int countList(){
    Node* currentNodePtr = headPtr->getNextNodePtr();
    int numElements = 0;

    while (currentNodePtr != 0){
        numElements++;
        currentNodePtr = currentNodePtr->getNextNodePtr();
    }
    return numElements;
}
```

```
void IterativePrint(){
    Node* currentNodePtr = headPtr->getNextNodePtr();

    while (currentNodePtr != 0){
        cout << currentNodePtr->getData() << " ";
        currentNodePtr = currentNodePtr->getNextNodePtr();
    }
    cout << endl;
}
```

C++

**C++**

```
bool containsElement(int searchData){
    Node* currentNodePtr = headPtr->getNextNodePtr();

    while (currentNodePtr != 0){

        if (currentNodePtr->getData() == searchData)
            return true;

        currentNodePtr = currentNodePtr->getNextNodePtr();
    }
    return false;
}
```

# Hash table Implementation (Code: 5.1)

```
private:                                Class Hashtable  
    List* listArray;                    C++  
    int tableSize;  
  
public:  
    Hashtable(int size){  
        tableSize = size;  
        listArray = new List[size];  
    }  
  
    int getTableSize(){  
        return tableSize;  
    }  
  
    void insert(int data){  
        int hashIndex = data % tableSize;  
        listArray[hashIndex].insert(data);  
    }  
}
```

```
void deleteElement(int data){  
    int hashIndex = data % tableSize;  
    while (listArray[hashIndex].deleteElement(data));  
}
```

**Class Hashtable**  
**C++**

To delete all the entries for the data  
in the Linked List

```
bool hasElement(int data){  
  
    int hashIndex = data % tableSize;  
    return listArray[hashIndex].containsElement(data);  
  
}
```

```
void printHashTable(){  
    for (int hashIndex = 0; hashIndex < tableSize; hashIndex++){  
        cout << "Hash Index: " << hashIndex << " : " ;  
        listArray[hashIndex].IterativePrint();  
    }  
}
```

```
void deleteElement(int data){  
    int hashIndex = data % tableSize;  
    //while (listArray[hashIndex].deleteElement(data));  
    listArray[hashIndex].deleteElement(data);  
}
```

Just to delete the first occurrence  
of the data in the Linked List

# Sample Applications of Hash table

- To test whether a test sequence is a permutation of a given sequence
- To print the unique elements in an array
- To find the union of two linked lists

# Permutation Check (Code 5.2)

- Given a sequence of integers (as a string)
  - Use string tokenizer to parse the string, extract the individual integers and store them in a hash table
- Given the test sequence of integers (also a string)
  - Use string tokenizer to parse the string and extract the individual integers
    - For each integer in the test sequence, check if it is in the hash table.
      - If so, delete it (just delete the first occurrence of the integer in the appropriate linked list of the hash table)
      - Otherwise, STOP and say the test sequence is not a permutation of the original sequence
    - If the hash table is empty after the test sequence is processed, print that the test sequence is a permutation of the original sequence; otherwise, it is not a permutation of the original.

# Hashtable: isEmpty( )

```
bool isEmpty(){
```

**Code 5.2: C++**

```
    for (int hashIndex = 0; hashIndex < tableSize; hashIndex++){
```

```
        if (!listArray[hashIndex].isEmpty())
```

```
            return false;
```

```
    }
```

**// If even one linked list is not empty, return false**

**// (i.e., the hash table is not empty)**

```
    return true;
```

**// If all the linked lists are empty, return true**

```
}
```

**// (i.e., the hash table is empty)**

```

string integerSequence; // Input the original integer sequence as a
cout << "Enter the integer sequence: "; // comma separated list
getline(cin, integerSequence); // Example: 45, 23, 12, 23, 90

string testSequence; // Input the test sequence likewise
cout << "Enter the test sequence for permutation: "; // Example: 12, 23, 45, 90, 23
getline(cin, testSequence);

int hashTableSize;
cout << "Enter the size of the hash table: "; // Input the size of the hash table and
cin >> hashTableSize; // Initialize a hash table of that size
Hashtable hashTable(hashTableSize);

char* integerArray = new char[integerSequence.length()+1];
strcpy(integerArray, integerSequence.c_str());

char* cptr = strtok(integerArray, ","); // Code for tokenizing the original
// integer sequence and extracting
// the individual integers

while (cptr != 0){
    string token(cptr);
    int value = stoi(token);

    hashTable.insert(value); // Insert the extracted integers in the hash table

    cptr = strtok(NULL, ",");
}

cout << endl;

```

**Code 5.2 (C++)**

```

hashTable.printHashTable(); // Print the contents of the hash table Code 5.2 (C++)

char* testArray = new char[testSequence.length()+1]; // Code to string tokenize the
strcpy(testArray, testSequence.c_str()); // test integer sequence and extract
char* tptr = strtok(testArray, ", "); // the individual integers

while (tptr != 0){
    string token(tptr); // Check whether a test integer is in the
    int testValue = stoi(token); // hash table; if so, delete it.
    if (hashTable.hasElement(testValue)) // If not in the hash table, print the test
        hashTable.deleteElement(testValue); // sequence is not a permuted seq. STOP!
    else{
        cout << testSequence << " is not a permuted sequence of " << integerSequence << endl;
        return 0; // equivalent to stopping the program right away!
    }

    tptr = strtok(NULL, ", ");
} // If the hash table is empty when we finish processing the test sequence, it implies
// there exists an element in the original sequence for every element in the test
// sequence and all of these elements were deleted. Hence, they are a permutation
if (hashTable.isEmpty())
    cout << testSequence << " is a permuted sequence of " << integerSequence << endl;
else
    cout << testSequence << " is not a permuted sequence of " << integerSequence << endl;

```

# Printing the Unique Elements (Code 5.3)

- Given an array  $A[0 \dots n-1]$  that may have elements appearing more than once, we could use the hash table to store the unique elements and print them.
- For every element  $A[i]$ , with  $0 \leq i \leq n-1$ , we store the element  $A[i]$  in the hash table the first time we come across it as well as print it.
- Hence, for every element  $A[i]$ , with  $0 \leq i \leq n-1$ , we check if  $A[i]$  is already in the hash table.
  - If  $A[i]$  is not already in the hash table, it implies  $A[i]$  has not been seen before: so, we print it out as well as insert in the hash table.
  - If  $A[i]$  is already in the hash table, it implies we have already printed it out and should not be printed again.
- The time complexity of the algorithm is dependent on the time to check whether a particular element is in the hash table or not for all values of the array index. If this could be done in  $\Theta(1)$  time per element, the asymptotic time complexity of the algorithm is  $\Theta(n)$ .

### Code 5.3 (C++)

```
int numElements;
cout << "Enter the number of elements you want to store in the array: ";
cin >> numElements;    // Input the number of elements to store in the array

int maxValue;
cout << "Enter the maximum value for an element: ";
cin >> maxValue;      // Input the maximum value for an element

int hashTableSize;
cout << "Enter the size of the hash table: ";
cin >> hashTableSize; // Input the number of indexes in the hash table

srand(time(NULL));    // Initialize the random number generator

int array[numElements]; // Generate the random elements
cout << "Elements generated: "; // and store them in the array as well as
for (int index = 0; index < numElements; index++){ // print them
    array[index] = rand() % maxValue;
    cout << array[index] << " ";
}

cout << endl;
```

```

Hashtable hashTable(hashTableSize); // Initialize the
// Hash table

for (int index = 0; index < numElements; index++){

    if (!hashTable.hasElement(array[index])){
        cout << array[index] << " ";
        hashTable.insert(array[index]);
    }

    // For an element array[index]: check if it is in the
    // hashTable. If it is not there, it implies the element
    // has not been seen before and print it as well as add
    // it to the hashTable.
    // If the element is already there in the hashTable,
    // it implies it is the duplicate entry for the element and
    // is not printed.
}

cout << endl;

```

**Code 5.3 (C++)**

# Find the Union of two Linked Lists (Code 5.4)

- Let L1 and L2 be the two linked lists and unionList be the union of the two lists
- unionList is initially empty and an element is to be included only once in this list (i.e., elements with duplicate entries in L1 or L2 are included only once in the unionList).
- We populate the contents of L1 in a hash table and unionList:
  - An element is added to both the hash table and unionList if the element is not in the hash table.
- We go through list L2, element by element. If an element in L2 is not in the hash table (it implies the element is not in L1 and unionList either), then it is included in the hash table as well as the unionList.

## Code 5.4 (C++)

```
int numElements;
cout << "Enter the number of elements you want to store in the two lists: ";
cin >> numElements;
    // Assume the two linked lists have the same number of elements
int maxValue; // The maximum value that could be for an element in the two lists
cout << "Enter the maximum value for an element: "; // is also the same
cin >> maxValue;

int hashTableSize;
cout << "Enter the size of the hash table: "; // Input the number of indexes for the
cin >> hashTableSize; // Hash table

srand(time(NULL)); // Initialize the random number generator

List firstList;
cout << "Elements generated for the first list: "; // generate random integers and
for (int index = 0; index < numElements; index++){ // populate the firstList
    int value = rand() % maxValue;
    firstList.insert(value);
}
firstList.IterativePrint();

List secondList;
cout << "Elements generated for the second list: "; // generate random integers and
for (int index = 0; index < numElements; index++){ // populate the secondList
    int value = rand() % maxValue;
    secondList.insert(value);
}
secondList.IterativePrint();
```

**Code 5.4 (C++)**

```
List unionList; // The unionList will have only unique elements
Hashtable hashTable(hashTableSize); // elements

for (int index = 0; index < numElements; index++){
    int value = firstList.read(index);

    if (!hashTable.hasElement(value)){ // Scan through the firstList and
        hashTable.insert(value); // insert an element in both the
        unionList.insert(value); // hash table and unionList if the
    } // element is not already in
    // the hash table
}

for (int index = 0; index < numElements; index++){
    int value = secondList.read(index); // Scan through the secondList
    // If an element in the secondList
    // is not in the hash table, it implies
    // the element is not in the unionList
    // either and hence add the element
    // to both the unionList and the
    // hash table

    if (!hashTable.hasElement(value)){
        hashTable.insert(value);
        unionList.insert(value);
    }
}

cout << "Elements in the union list: ";
unionList.IterativePrint();
```