

# Small World Networks

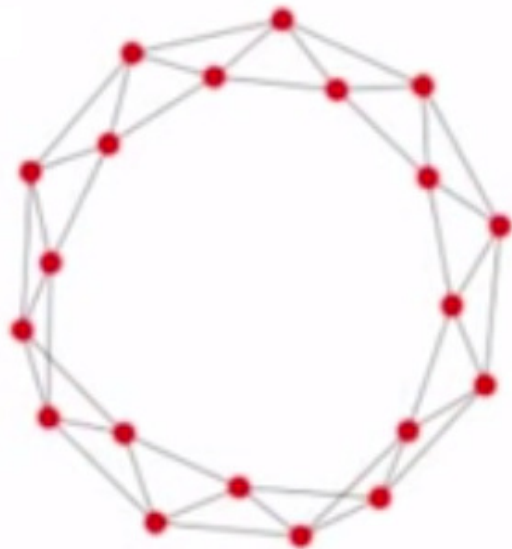
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# Small-World Networks

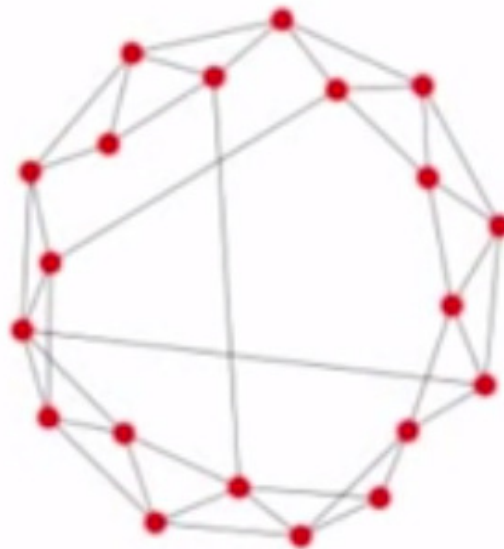
- A small-world network is a type of graph in which most nodes are not neighbors of one another, but most nodes can be reached from every other by a small number of hops.
- Specifically, a small-world network is defined to be a network where the typical distance  $L$  (the number of hops) between two randomly chosen nodes grows proportionally to the logarithm of the number of nodes in the network.
- Examples of Small-World Networks:
  - Road maps, food chains, electric power grids, metabolite processing networks, networks of brain neurons, voter networks, telephone call graphs, gene regulatory networks.

# Small Worlds

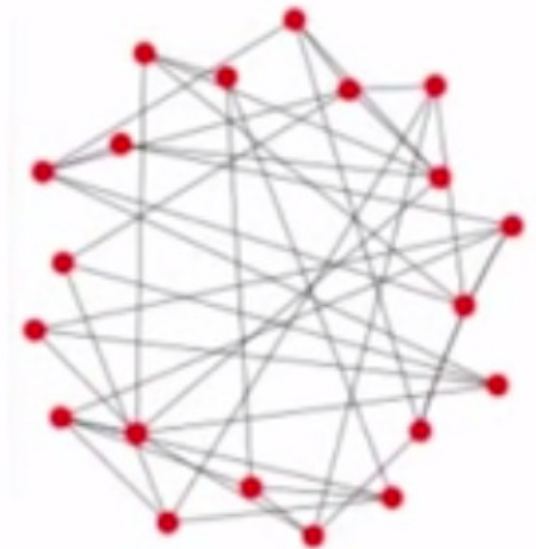
- Two major properties of small world networks
  - High average clustering coefficient
    - The neighbors of a node are connected to each other
    - Nodes' contacts in a social network tend to know each other.
  - Short average shortest path length
    - Shorter paths between any two nodes in the network



(Regular graph)



(Small-world network)



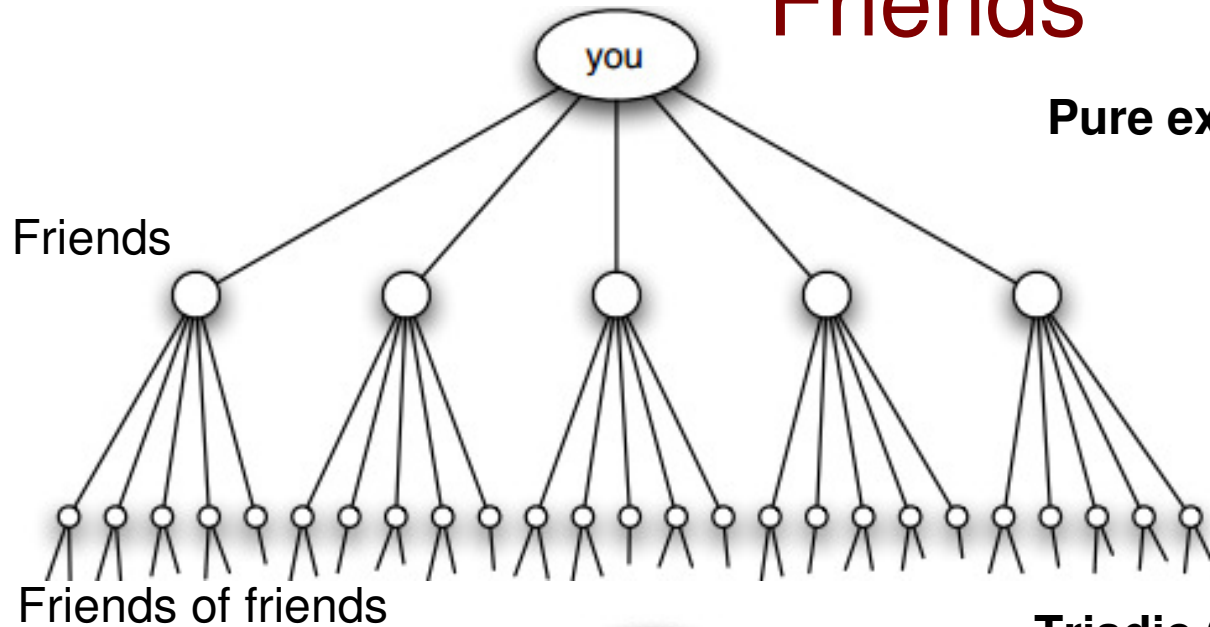
(Random graph)

0

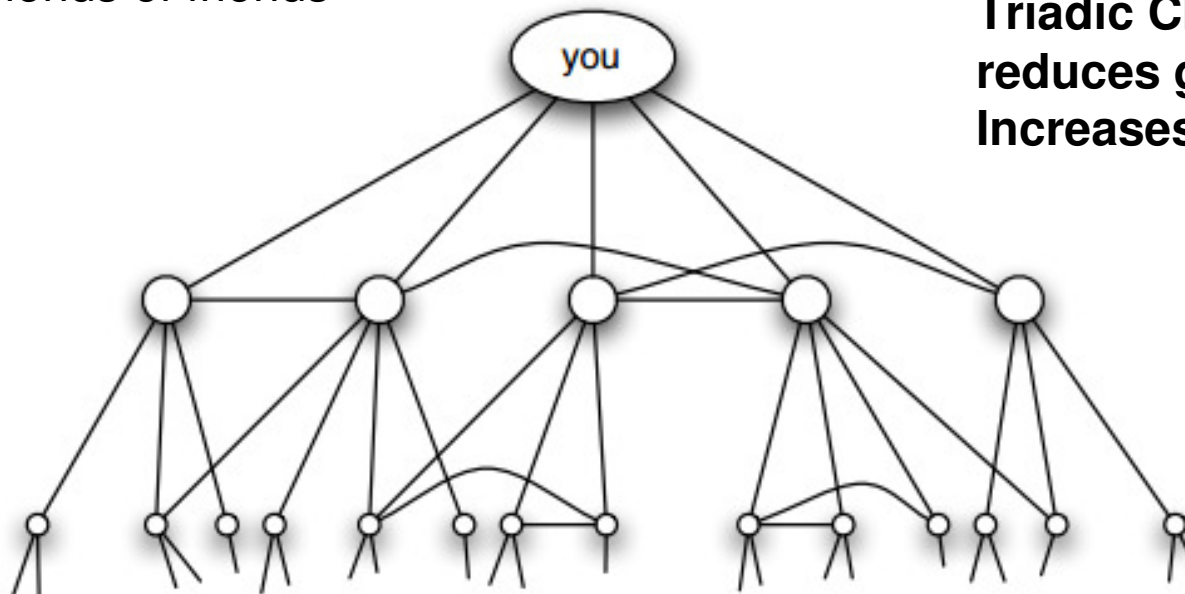
Randomness

1

# Triadic Closure and Growth Rate of New Friends



**Pure exponential growth**  
**No triadic closure**  
**Zero clustering coefficient**

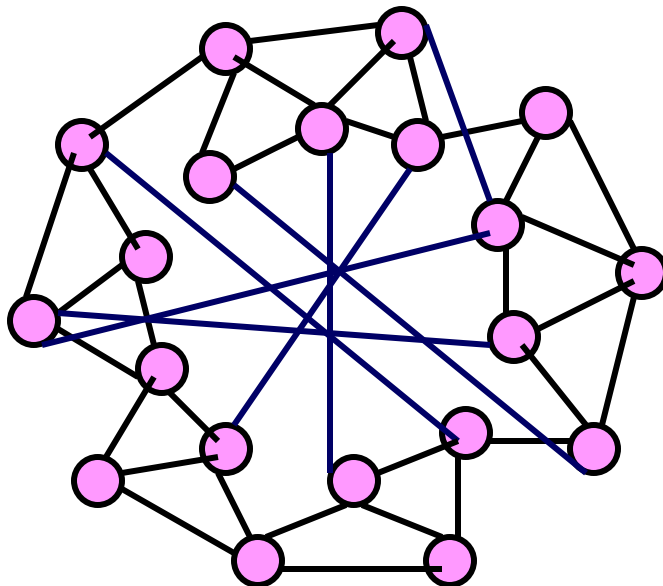
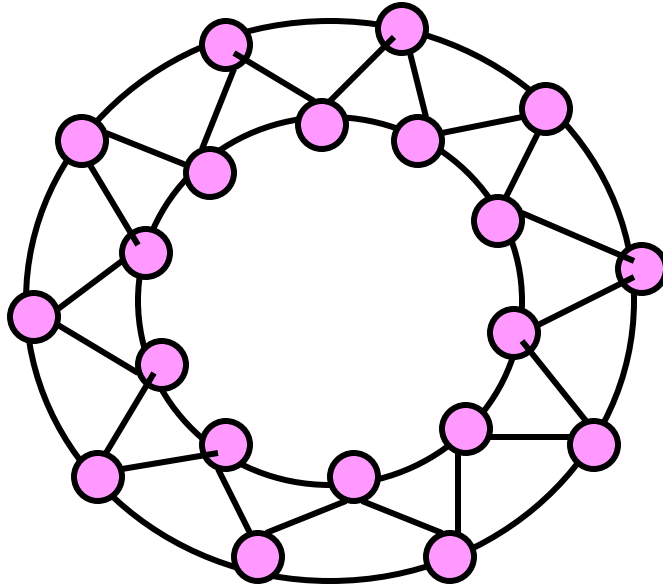


**Triadic Closure**  
**reduces growth rate of new friends**  
**Increases clustering coefficient**

# Modeling Small World Networks

- The ER model for random graphs provided shorter paths between any two nodes in the network. However, the ER graphs have a low clustering coefficient and triadic closures.
  - ER graphs have a constant, random and independent probability of two nodes being connected.
- The Watts and Strogatz model (WS model) accounts for clustering while retaining the short average path lengths of the ER model.
- The WS model interpolates between an ER graph and a regular ring lattice.

# WS Model



- Watts and Strogatz (WS) Model:  
The WS model interpolates between an ER graph and a regular ring lattice.

- Let  $N$  be the number of nodes and  $K$  (assumed to be even) be the mean degree.
- Assume  $N \gg K \gg \ln(N) \gg 1$ .
- There is a rewiring parameter  $\beta$  ( $0 \leq \beta \leq 1$ ).
- Initially, let there be a regular ring lattice of  $N$  nodes, with  $K$  neighbors ( $K/2$  neighbors on each side).
- For every node  $n_i = n_0, n_1, \dots, n_{N-1}$ , rewire the edge  $(n_i, n_j)$ , where  $i < j$ , with probability  $\beta$ . Rewiring is done by replacing  $(n_i, n_j)$  with  $(n_i, n_k)$  where  $n_k$  is chosen uniformly-randomly among all possible nodes that avoid self-looping and link duplication.

$\beta = 0 \rightarrow$  Regular ring lattice

$\beta = 1 \rightarrow$  Random network

# Small-World Network: WS Model

- The underlying lattice structure of the model produces a locally clustered network, and the random links dramatically reduce the average path lengths
- The algorithm introduces about  $(\beta NK/2)$  non-lattice edges.
- Average Path Length ( $\beta$ ):
  - Ring lattice  $L(0) = (N/2K) \gg 1$
  - Random graph  $L(1) = (\ln N / \ln K)$
  - For  $0 < \beta < 1$ , the average path length reduces significantly even for smaller values of  $\beta$ .

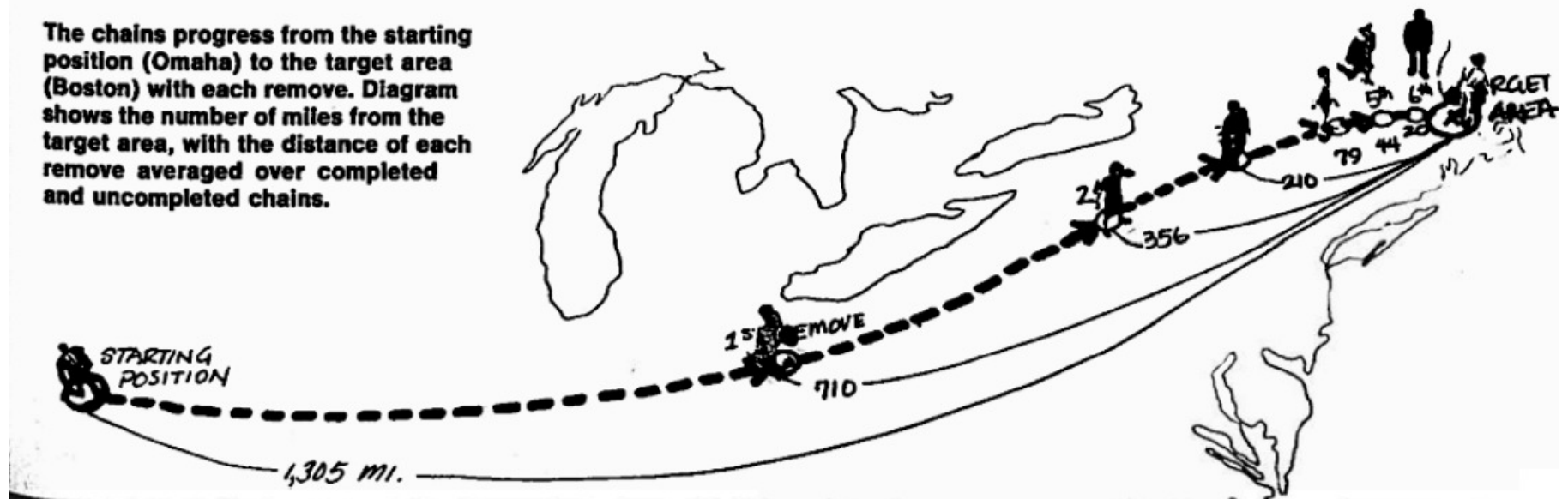
- Clustering Coefficient ( $\beta$ ):

$$C(0) = \frac{3(K-2)}{4(K-1)} \quad C(1) = \frac{K}{N} \quad C'(\beta) = C(0) * (1 - \beta)^3$$

- For  $0 < \beta < 1$ , the clustering coefficient remains close to that of the regular lattice for low and moderate values of  $\beta$  and falls only at relatively high  $\beta$ .
- For low-moderate values of  $\beta$ , we thus capture the small-world phenomenon where the average path length falls rapidly, while the clustering coefficient remains fairly high.

# Limitations of the WS Model

- The WS model introduced the notion of random edges to infuse shorter path lengths amidst larger clustering coefficient.
- However, the long-range edges span between any two nodes in the network and do not mimic the edges of different lengths seen in real-world networks (like in the US road map as in Milgram's experiment or airline map).
  - Path lengths could not be as small as they are in real networks.
  - Need some edges to nodes that are few hops away, rather than edges to some arbitrarily chosen nodes.
  - Cannot generate hubs as in scale-free networks.



Source: Figure 20.4: Easley and Kleinberg



# Enhancement to the WS Model

- In addition to the re-wiring parameter  $\beta$ , another parameter called the clustering exponent ( $q$ ) is introduced.
- An  $(u, v)$  edge is selected for re-wiring with a probability  $\beta$ . After being selected, we do not randomly re-wire  $u$  with a node  $w$ . Instead, we pick a pair  $(u, w)$  for re-wiring with a probability of  $[d(u, w)^{-q}] / 2\log n$ , where
  - For optimal results,  $q$  must be the dimensionality of the network modeled. **For a ring lattice,  $q = 1$ .**
  - $n$  is the number of nodes in the network.
  - $d(u, w)$  is the minimum number of hops between  $u$  and  $w$  in the original network layout (before enhancement)
    - The ring lattice is a single-dimension network
    - A grid is a two-dimensional network.
  - To implement this enhancement in simulations, we generate a random number between 0 to 1; the  $(u, w)$  pair whose  $[d(u, w)^{-dim}] / 2\log n$  value is closest and above the random number generated is chosen for re-wiring.
- With this re-wiring model, if routed optimally, (on average) the # hops in the path to the target is expected to reduce by a factor of 2 with every additional hop in the path  $[\log n \text{ hops}]$